

# An ultrafast quantum random number generator based on quantum phase fluctuations

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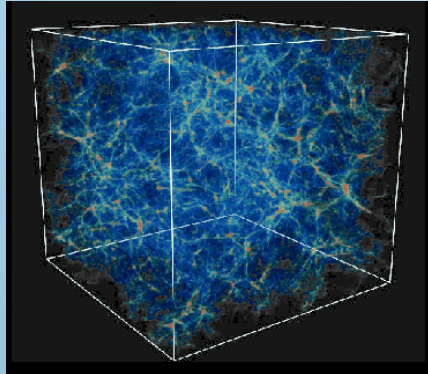
QCRYPT2011, ETH, Zurich,  
Sep. 14<sup>th</sup>, 2011

# Outline

- I** Introduction
- II** Experimental setup and results
- III** Post-processing
- IV** Future directions

# Applications of random numbers

Scientific simulations



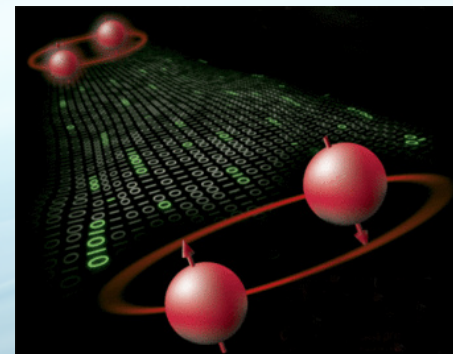
Lottery & Gambling



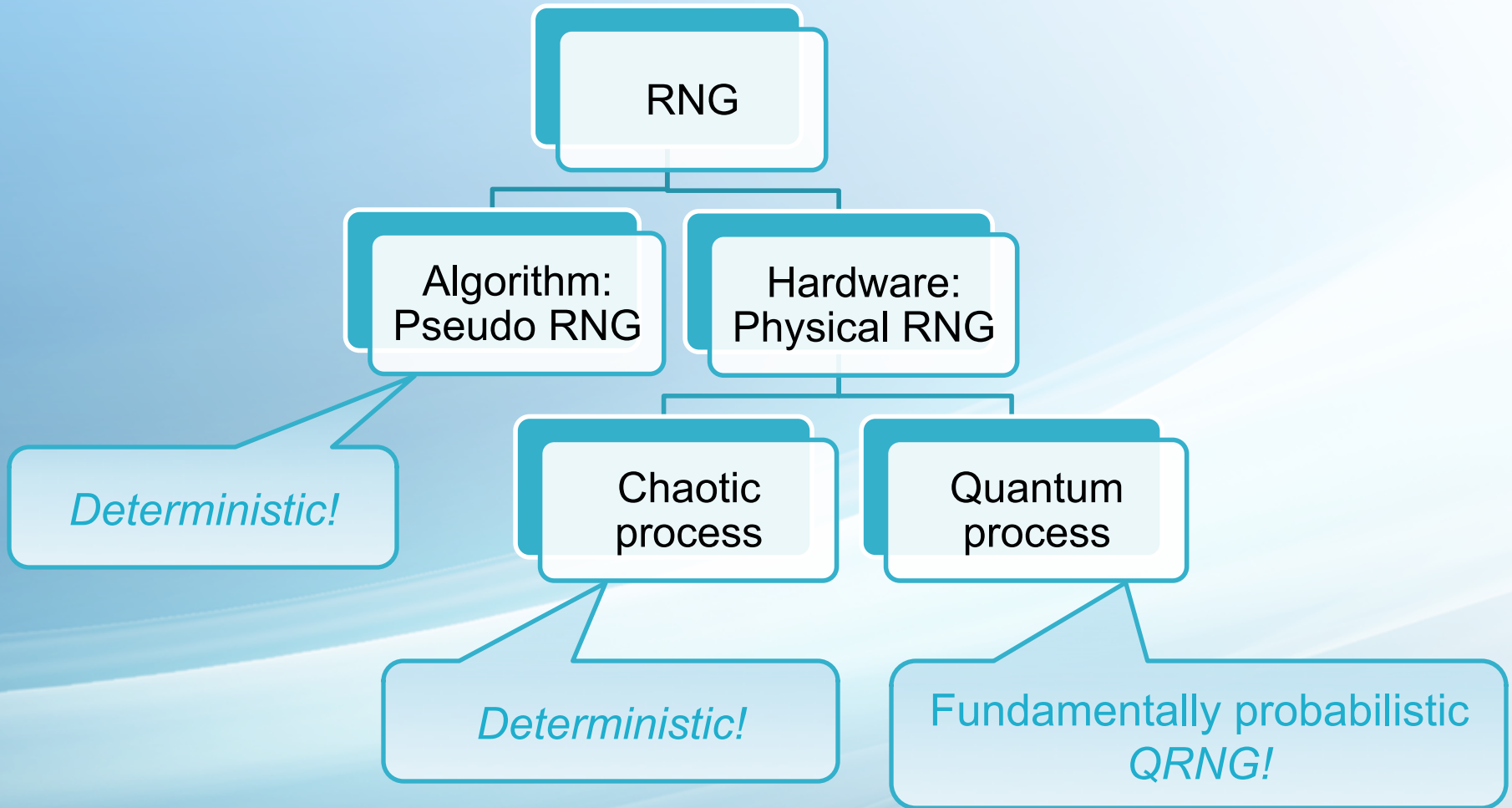
Market



Cryptography

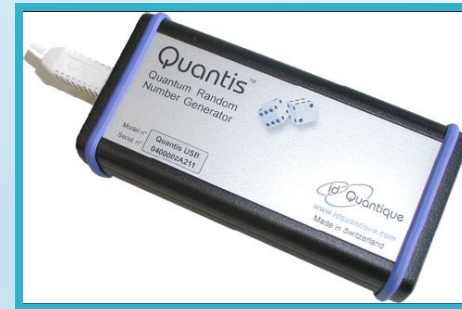
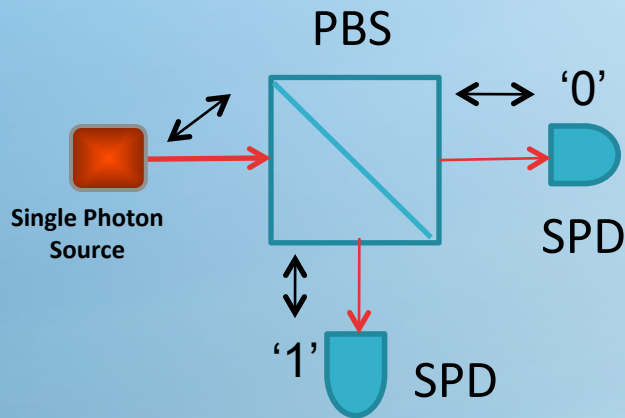


# Random number generator (RNG)



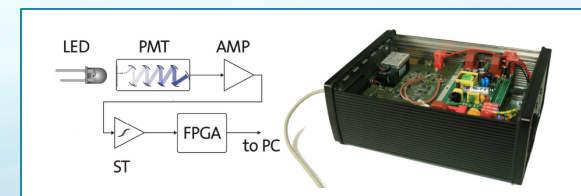
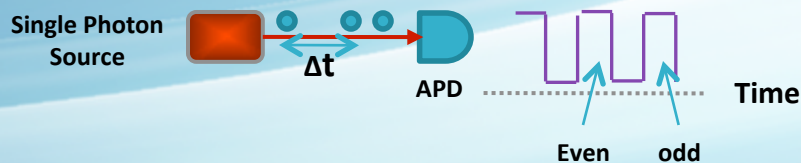
# QRNG: single photon detection

- Polarization measurement [1]



Commercial QRNGs up to 16 Mb/s.  
(Figure is from ID Quantique)

- Photon arrival time [2-3] or photon number counting [4]

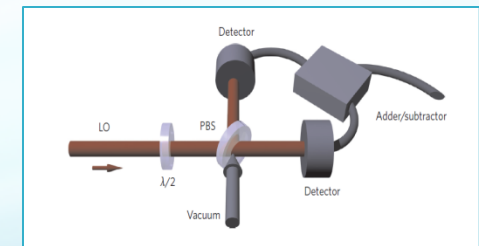
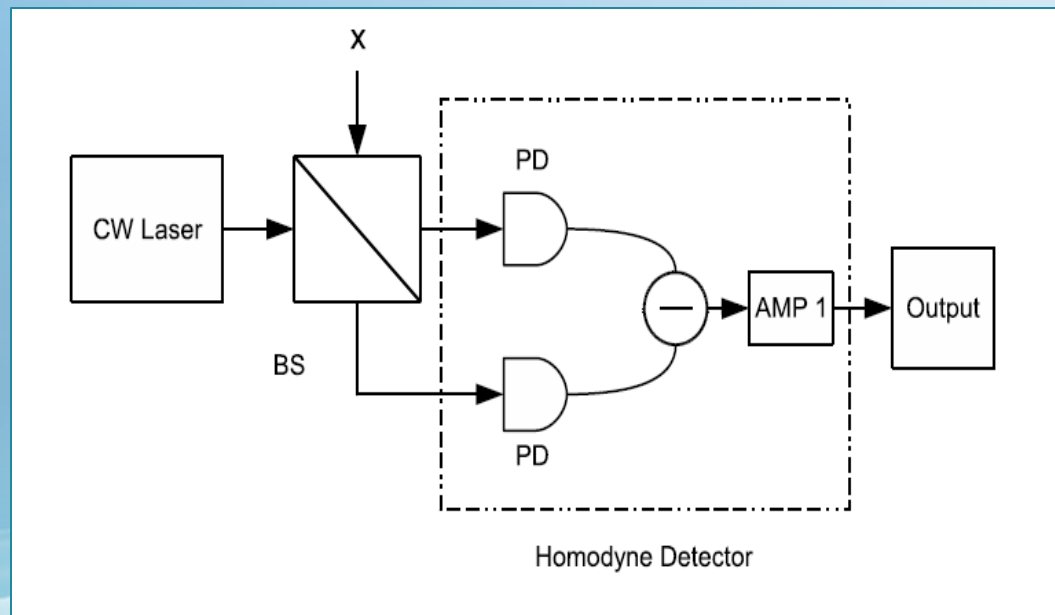


Fully integrated QRNG.  
(Figure is from ref [4])

- [1] T. Jennewein, et al, Rev. of Sci. Ins., 71:1675-1680, 2000.
- [2] P. Kwiat, E. Jeffrey, P. Altepeter, US Patent Appl. 20060010182, 2006.
- [3] J. Dynes, et al, App. Phys. Lett., 93, 031109 (2008)
- [4] M. Furst, et al, Opt. Exp. 18, 13029 (2010).

# QRNG: vacuum state fluctuations [1-2]

- Homodyne detection measuring the electrical field fluctuations of Vacuum state.



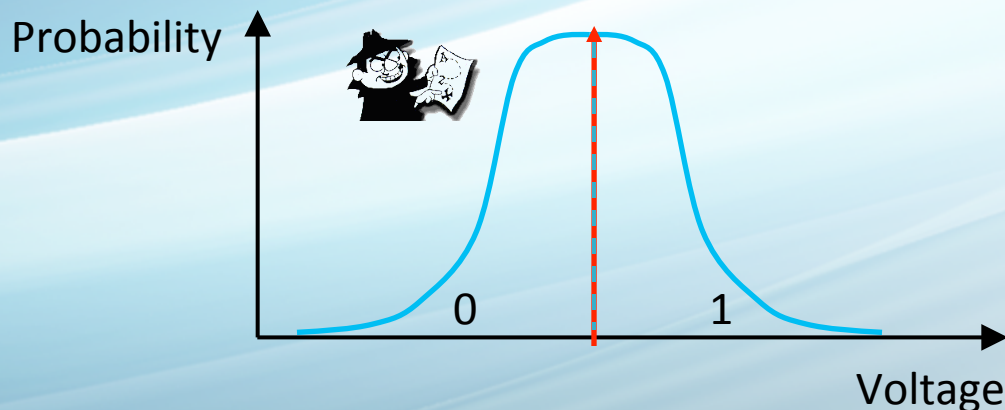
QRNG with 6.5 Mb/s [2].  
(Figure is from ref [2])

[1] A. Trifonov and H. Vig, US Patent No. 7,284,024, 16 October 2007

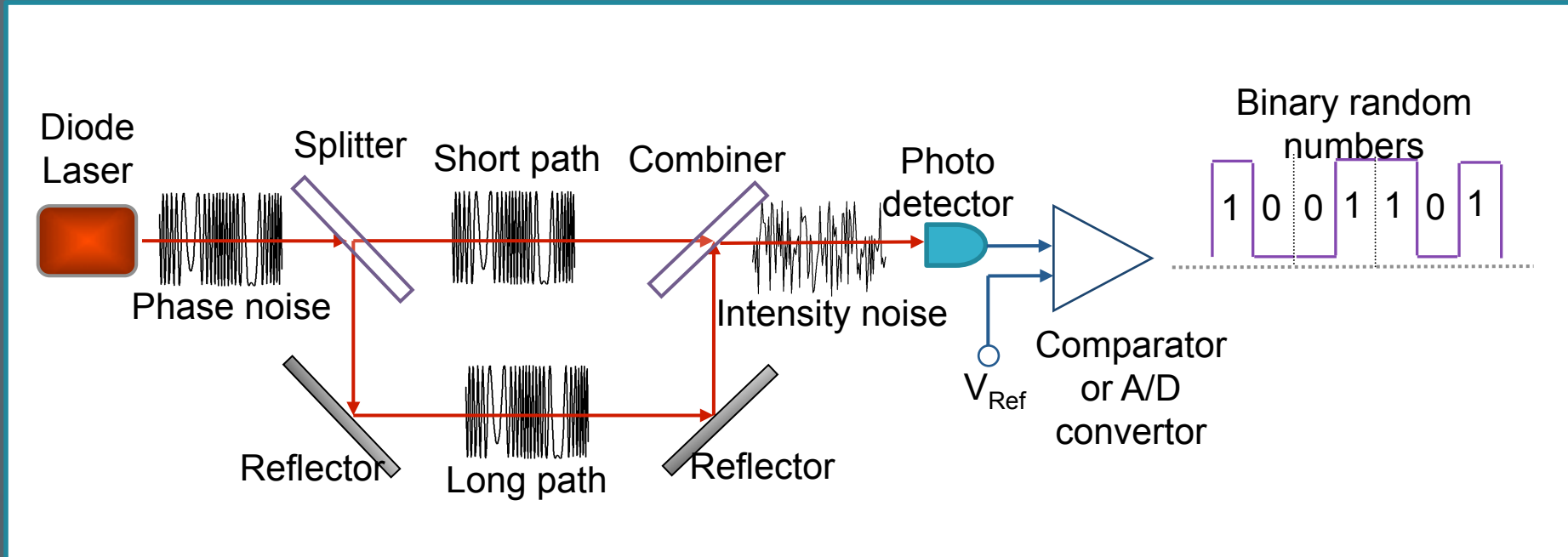
[2] C. Gabriel, et al, Nature Photonics, 4, 711–715 (2010)

# Motivation: QRNG existing problems

- Low generation rate
  - Typical rates:  $6.5 \text{ Mb/s}$  using vacuum state fluctuations,  $16 \text{ Mb/s}$  using polarization measurement (commercial QRNG),  $152 \text{ Mb/s}$  using photon arrival time [M. Wahl, et al, APL, 98, 171105, 2011].
- High cost
  - For example, the IDQ system (Quantis,  $16 \text{ Mb/s}$ ) costs 2230 €.
- Eavesdropper (Eve) may have partial information
  - Control side information (detector noise, environmental noise, etc).



# Our approach: randomness from laser phase fluctuations



Measure laser  
phase fluctua

Sample

Generate  
digital bits

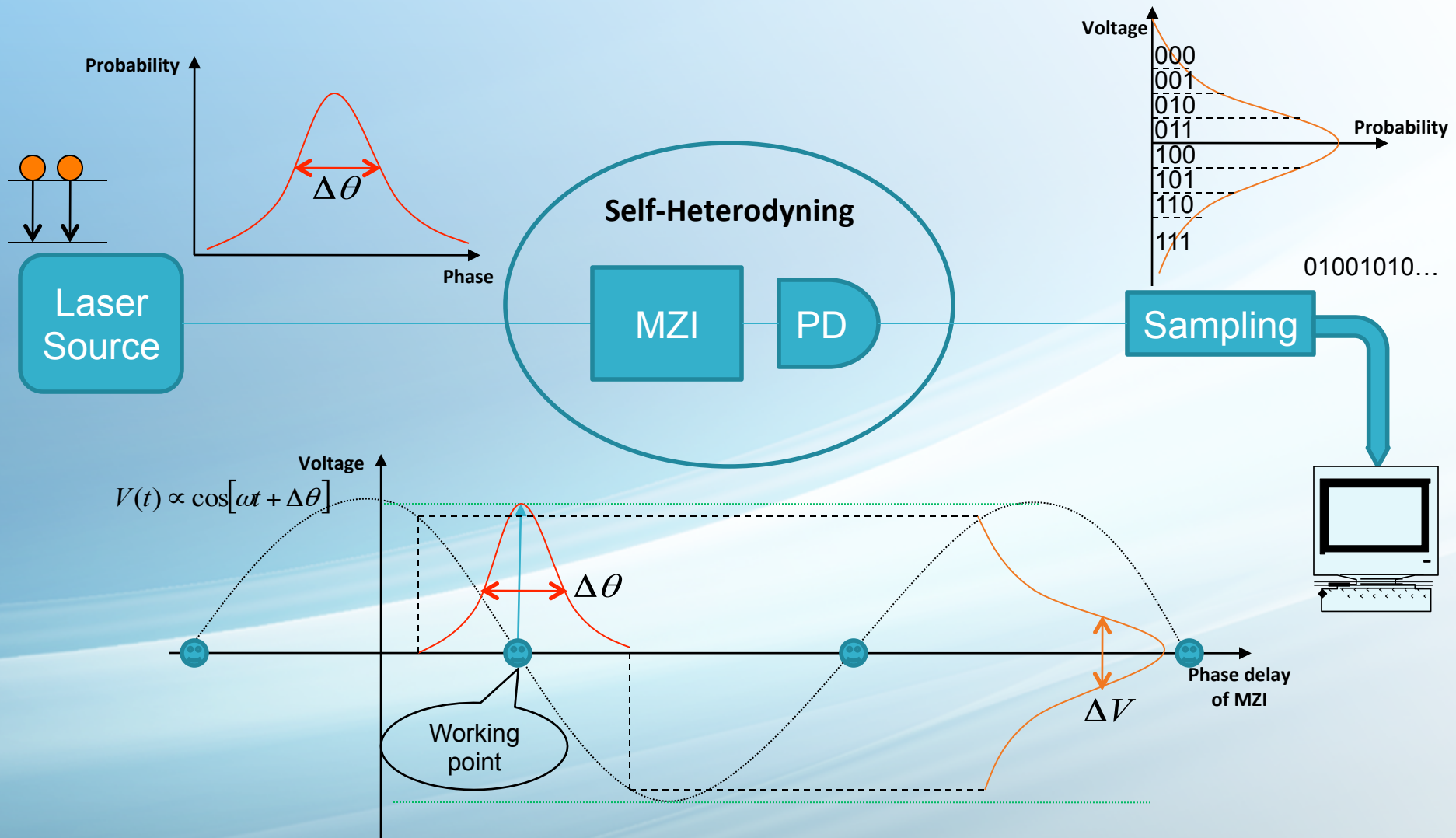
**Low cost & high rate!**

[1] A. Yariv and P. Yeh, "Photonics: optical electronics in modern communications" (6th edition), Oxford University Press (2007).

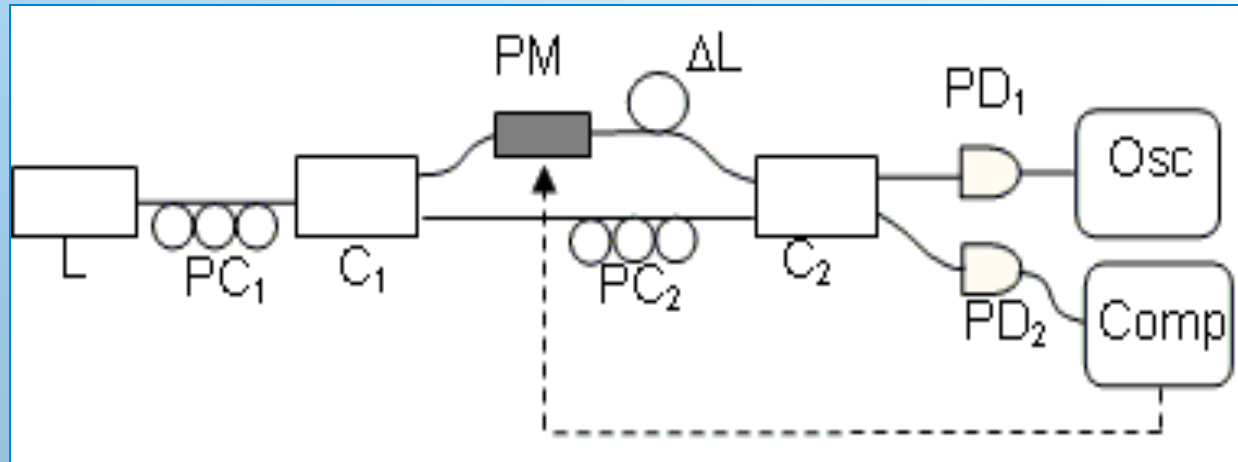
[2] K. Petermann, "Laser diode modulation and noise", (Springer, 1988).



# How the system works?



# Our previous work



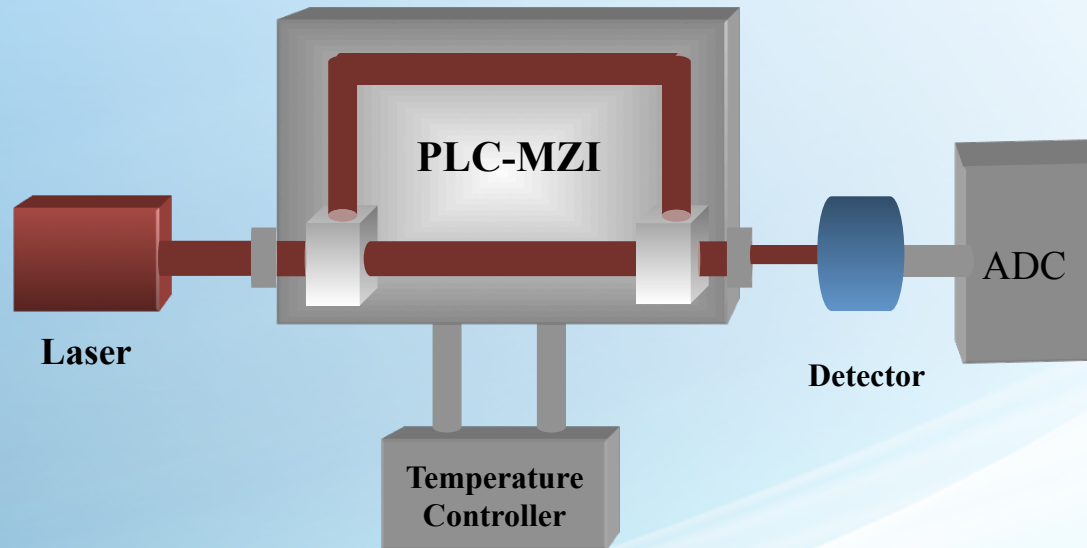
L: 1550nm cw DFB laser diode;  
PC1,2: polarization controller;  
PD2: 1MHz photo-receiver;  
OSC: 3GHz oscilloscope;

$C_{1,2}$ : fiber couplers  
PD1: 5GHz photo-detector;  
PM: phase modulator;  
Comp: computer with DAQ.

- Self-heterodyne system with off-the-shelf components.

- 500 Mb/s

# Our new setup



**PLC-MZI:** planar lightwave circuit Mach-Zehnder interferometer;

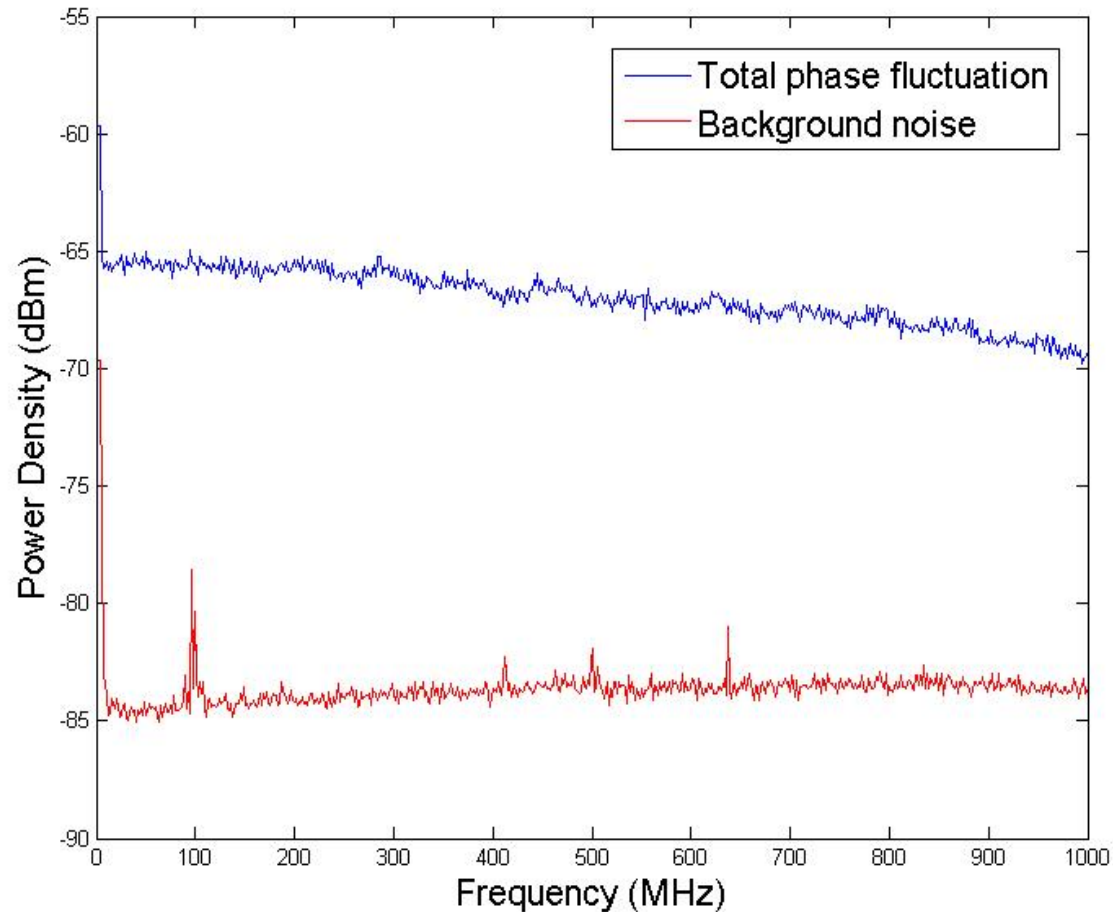
**ADC:** 8-bit analog-to-digital converter.

**Sampling rate:** 1G samples per second

**Extractable random bits:** 6.7 bits/ sample

Generation rate over **6 Gb/s!**

# Measurement results



Quantum phase fluctuation is dominant!

F. Xu, B. Qi, X. Ma, H. Xu, H. Zheng, H.-K. Lo, *arXiv: 1109.0643* (2011)

# Quantum signal and classical noise

- Laser phase fluctuations [1]
  - Quantum: spontaneous emission — Inversely power-dependent (Q/P)
  - Classical: cavity instability, etc. — power-independent (C)
- Electrical noise (detector) and EM noise (environment) — F
- Quantify the parameters:

$V(t) \propto c$

$$\langle V^2 \rangle = AP^2 \langle \Delta\theta^2 \rangle + F$$

$$= AP^2 \left( \frac{Q}{P} + C \right) + F = AQP + ACP^2 + F$$

delay  
|Z|

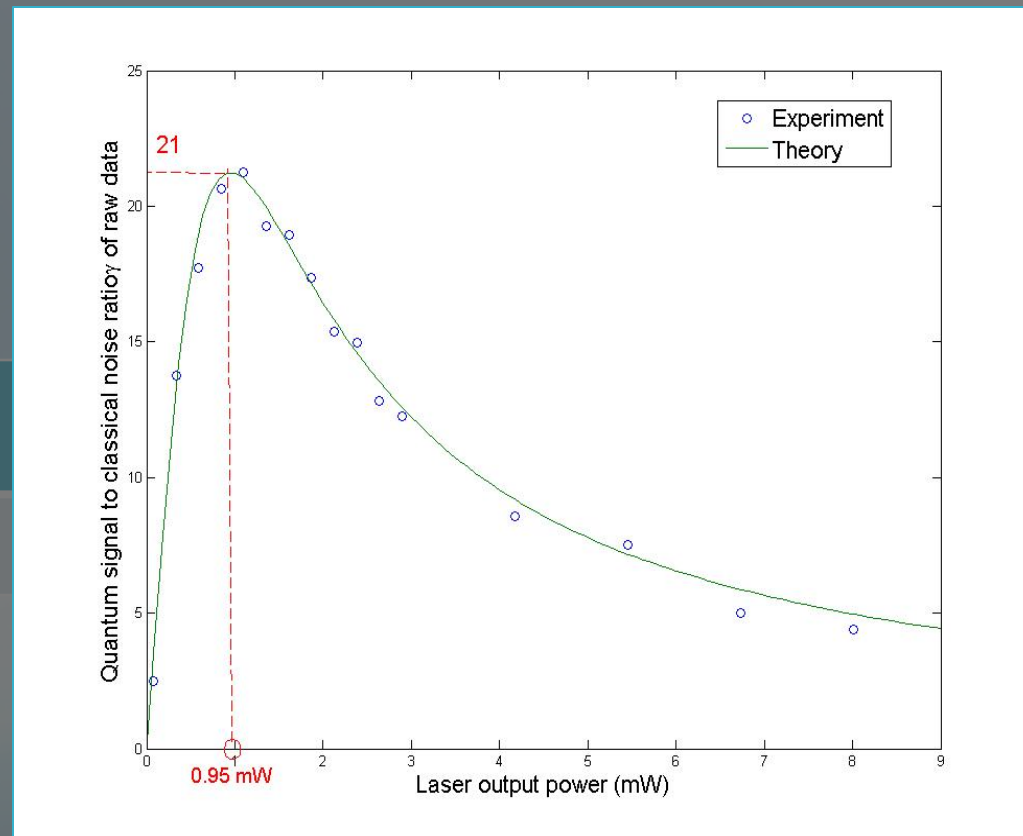
[1] C. H. Henry, IEEE J. Quantum Electron. QE-18, 259 (1982).

# Quantum signal and classical noise

$$\langle V^2 \rangle = AQP + ACP^2 + F$$

F (mV <sup>2</sup> )	AQ (mV <sup>2</sup> /mW)	AC (mV <sup>2</sup> /mW <sup>2</sup> )
0.36±0.06	16.12±0.49	0.40±0.16

$$\gamma = \frac{AQP}{ACP^2 + F}$$



# Post-processing

- Why we need post-processing?
  - Eve may have partial information (by controlling classical noise).
  - Quantum fluctuation is a non-uniform distribution (Gaussian).
- Extract out *uniform-quantum* randomness!

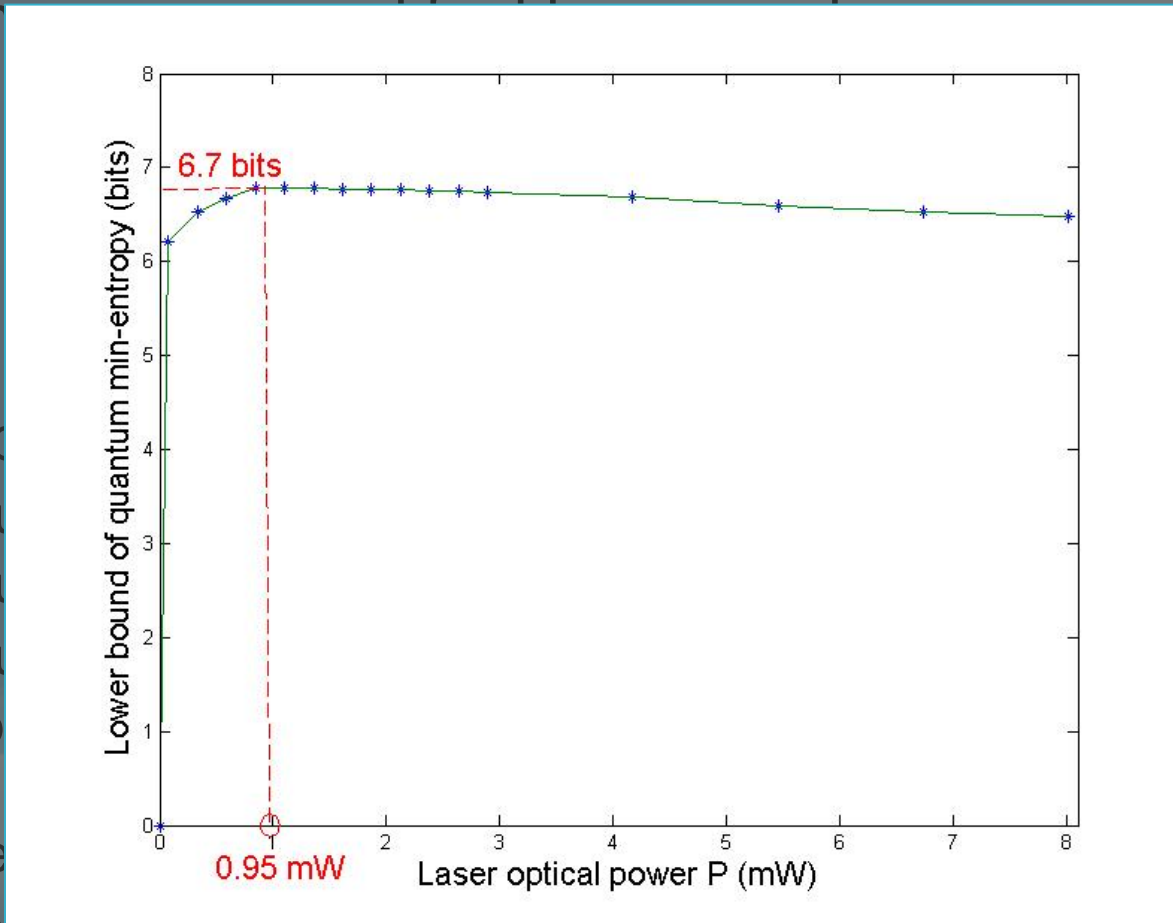
## Randomness extractor!

- Procedure
  - Min-entropy evaluation
  - Randomness extraction

# Min-entropy evaluation

- Random

- Quantum
- Assump
  - Quantum
  - Quantum
  - Quantum
  - Total p noise,
  - The se distribu



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# Randomness extraction

- Implement two extractors
  - Universal Hashing [1]
    - With Toeplitz matrix
  - Trevisan's Extractor [2]

Details of implementations:

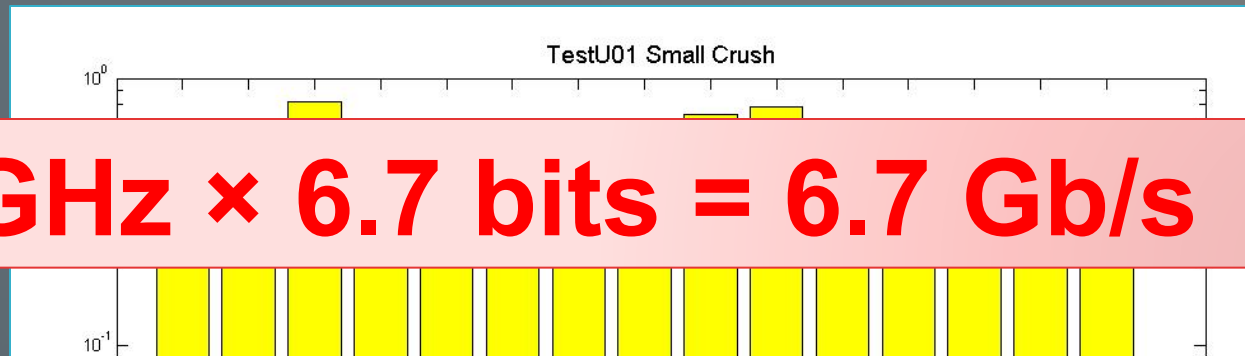
Xiongfeng Ma, et al, *under preparation* (2011)

- **QRNG with information-theoretically proven randomness!**

[1] M. Wegman and J. Carter, Journal of computer and system sciences 22, 265 (1981).

[2] L. Trevisan, Journal of the ACM 48, 2001 (1999).

# Extraction results: universal hashing



**1 GHz × 6.7 bits = 6.7 Gb/s**

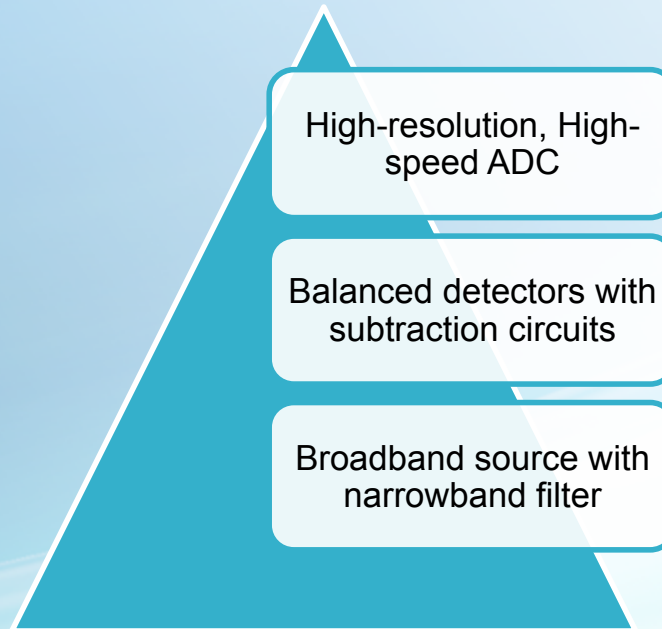
## Summary

- Demonstrate a simple and fast QRNG over 6 Gb/s!
- Quantify the quantum randomness by min-entropy!
- Implement two randomness extractors to extract out the quantum randomness!



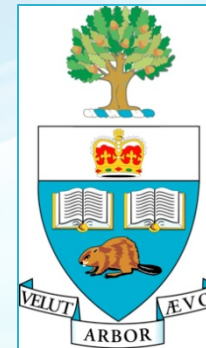
# Future directions

- Optimize system design.



- High-speed electronics for real time randomness extraction.
- Random number storage & transfer.

# Acknowledgements



Previous works:

- B. Qi, et al, *AQIS* (2009)
- B. Qi, et al, *Opt. Lett.* 35, 312 (2010)

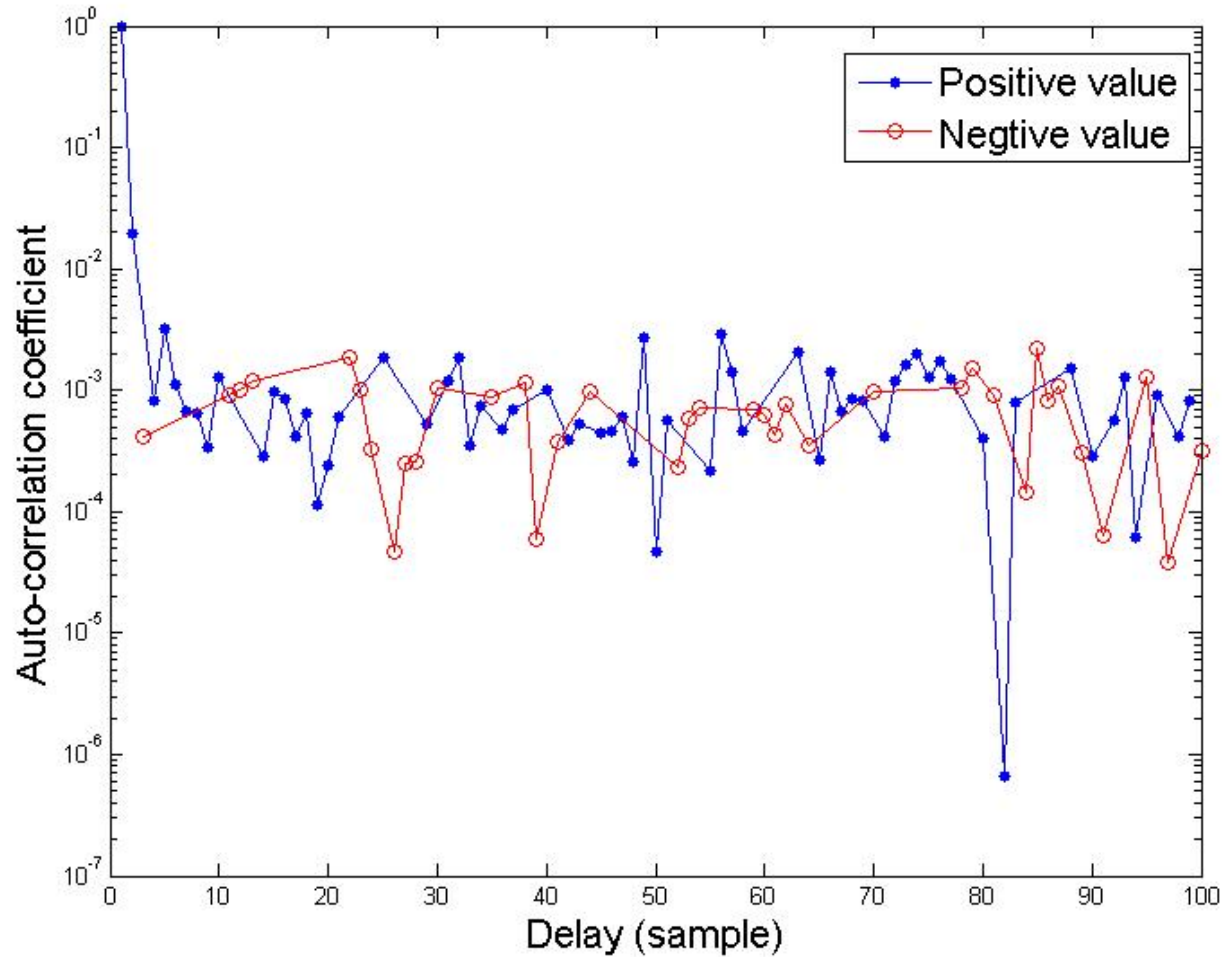
Current works:

- Feihu Xu, et al, *arXiv: 1109.0643* (2011)
- Xiongfeng Ma, et al, *under preparation* (2011)

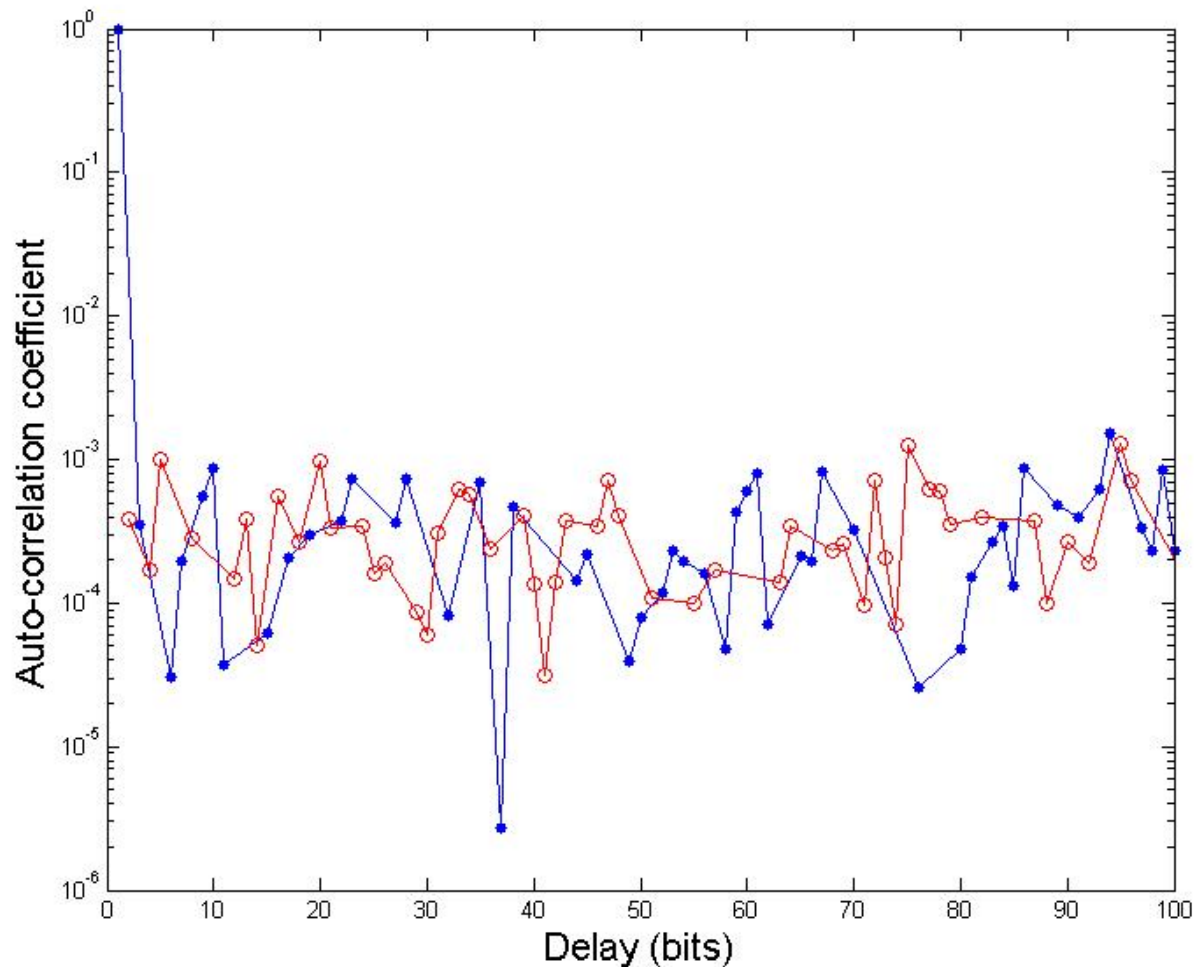
# Thank You !

# Backup part

# Autocorrelation of raw data

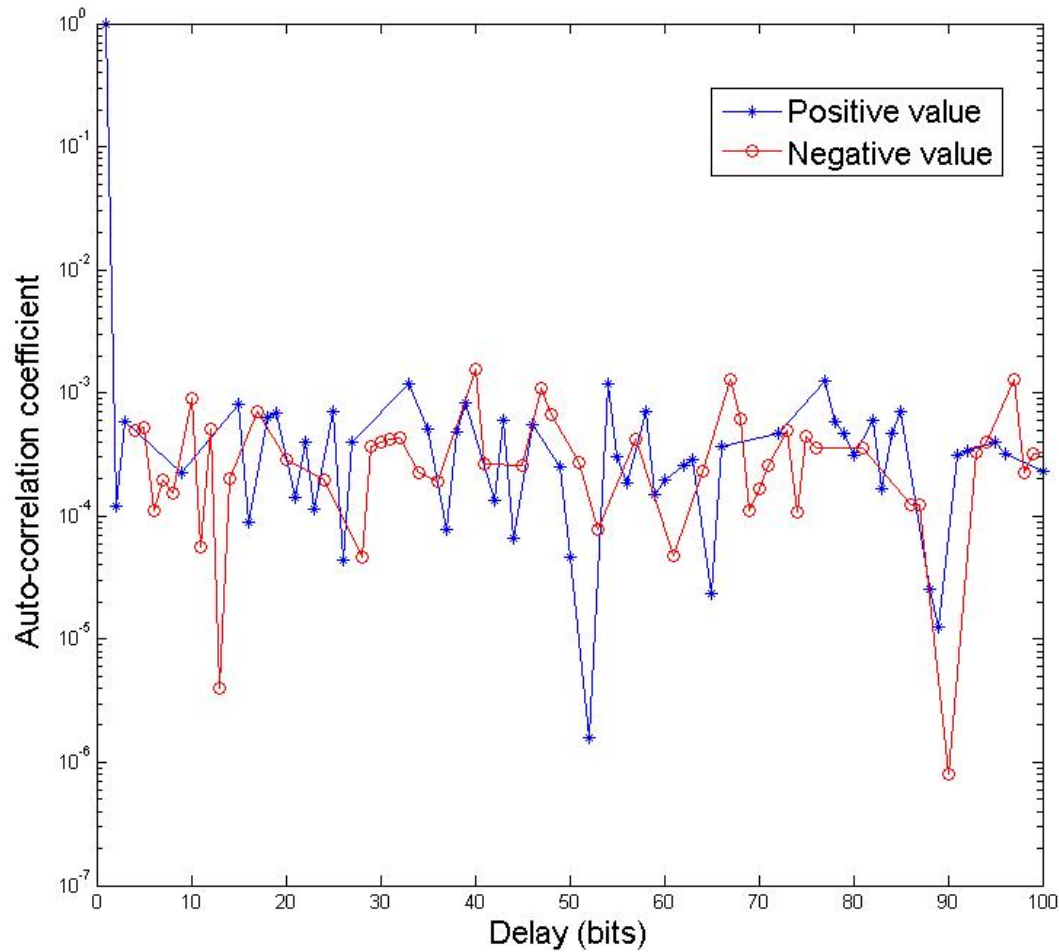


# Autocorrelation of Toeplitz-hashing output





# Autocorrelation of Trevisan's extractor output



# Diehard

	Pseudo-RNG	Raw data	Trevisan's		Toeplitz-hashing	
Statistical test	Result	Result	p-value	result	p-value	result
Birthday Spacings [KS]	success	<i>failure</i>	0.82263	success	0.340863	success
Overlapping permutations	success	<i>failure</i>	0.679927	success	0.403824	success
Ranks of 31x31 matrices	success	<i>failure</i>	0.419095	success	0.349441	success
Ranks of 31x32 matrices	success	<i>failure</i>	0.715705	success	0.816752	success
Ranks of 6x8 matrices [KS]	success	<i>failure</i>	0.195485	success	0.408573	success
Bit stream test	success	<i>failure</i>	0.048260	success	0.281680	success
Monkey test OPSO	success	<i>failure</i>	0.027300	success	0.892600	success
Monkey test OQSO	success	<i>failure</i>	0.023200	success	0.267200	success
Monkey test DNA	<i>failure</i>	<i>failure</i>	0.038000	success	0.736700	success
Count 1's in stream of bytes	success	<i>failure</i>	0.380162	success	0.639691	success
Count 1's in specific bytes	<i>failure</i>	<i>failure</i>	0.020417	success	0.373149	success
Parking lot test [KS]	<i>failure</i>	<i>failure</i>	0.629013	success	0.151689	success
Minimum distance test [KS]	success	<i>failure</i>	0.019499	success	0.688780	success
Random spheres test [KS]	success	<i>failure</i>	0.488703	success	0.939227	success
Squeeze test	success	<i>failure</i>	0.238004	success	0.155403	success
Overlapping sums test [KS]	success	<i>failure</i>	0.022339	success	0.909675	success
Runs test (up) [KS]	<i>failure</i>	<i>failure</i>	0.403504	success	0.181024	success
Runs test (down) [KS]	success	<i>failure</i>	0.119132	success	0.668512	success
Craps test No. of wins	success	<i>failure</i>	0.757521	success	0.826358	success
Craps test throws/game	success	<i>failure</i>	0.179705	success	0.862986	success

TABLE V: Diehard. Data size is 240Mbits. For the cases of multiple P-values, a Kolmogorov-smirnov (KS) test is used to obtain a final P-value, which measures the uniformity of the multiple P-values. The test is successful if all final P-values satisfy  $0.01 \leq P \leq 0.99$

	Pseudo-RNG	Raw data	Toeplitz-hashing		
Statistical test	Result	Result	p-value	Proportion	Result
Frequency	success	<i>failure</i>	0.373625	0.9900	success
Block-frequency	success	<i>failure</i>	0.310049	0.9960	success
Cumulative sums	success	<i>failure</i>	0.422638	0.9980	success
Runs	success	<i>failure</i>	0.703417	0.9900	success
LongestRun	success	<i>failure</i>	0.013569	0.9880	success
Rank	success	<i>failure</i>	0.411840	0.9940	success
FFT	success	<i>failure</i>	0.987079	0.9860	success
NonOverlappingTemplate	<i>failure</i>	<i>failure</i>	0.727851	0.9820	success
overlappingTemplate	success	<i>failure</i>	0.110083	0.9780	success
Universal	success	<i>failure</i>	0.962688	0.9880	success
ApproximateEntropy	success	<i>failure</i>	0.674543	0.9920	success
Random-excursions	success	<i>failure</i>	0.409207	0.9900	success
Random-excursions Variant	success	<i>failure</i>	0.426358	0.9840	success
Serial	success	<i>failure</i>	0.217570	0.9860	success
Linear-complexity	success	<i>failure</i>	0.657833	0.9940	success

TABLE VI: NIST. Data size is 3.25 Gbits (500 sequences with each sequence around 6.5 Mbits). To pass the test, P-value should be larger than the lowest significant level  $\alpha = 0.01$ , and the proportion of sequences satisfying  $P > \alpha$  should be greater than 0.976. Where the test has multiple P-values, the worst case is selected.

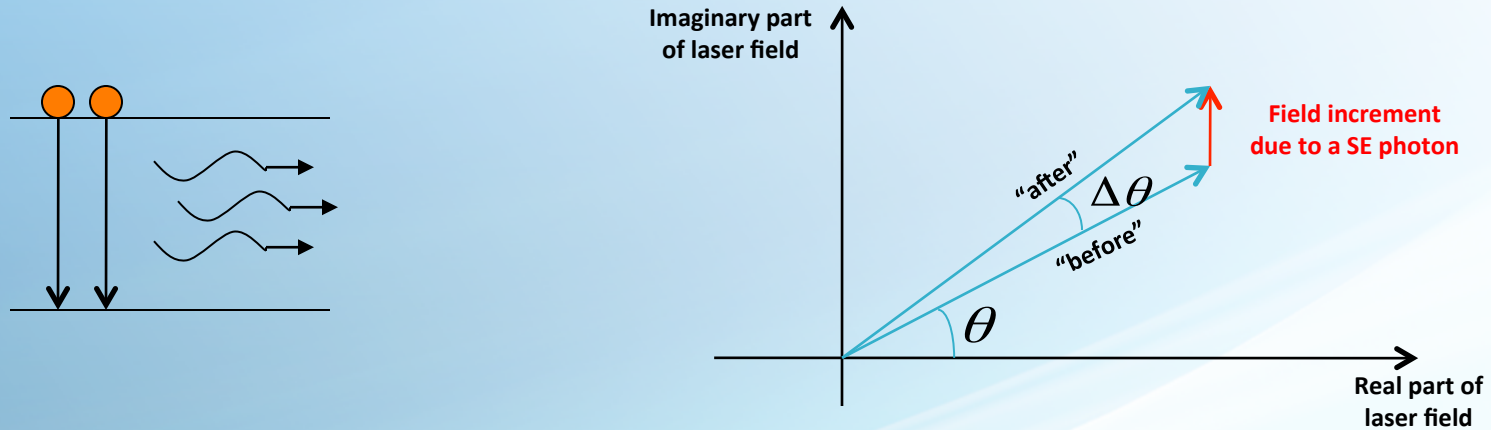
# TestU01

	Pseudo-RNG	Raw data	Toeplitz-hashing	
Statistical Test	Result	Result	p-value	Result
BirthdaySpacings	Success	<i>failure</i>	0.5300	success
Collision	Success	<i>failure</i>	0.1500	success
Gap Chi-square	success	<i>failure</i>	0.8900	success
SimpPoker Chi-square	success	<i>failure</i>	0.3500	success
CouponCollector Chi-square	success	<i>failure</i>	0.6700	success
MaxOft Chi-square	success	<i>failure</i>	0.6900	success
MaxOft Anderson-Darling	success	<i>failure</i>	0.9500	success
WeightDistrib Chi-square	success	<i>failure</i>	0.5600	success
MatrixRank Chi-square	success	<i>failure</i>	0.5100	success
Hammingindep Chi-square	success	<i>failure</i>	0.1000	success
RandomWalk1 H Chi-square	success	<i>failure</i>	0.9931	success
RandomWalk1 M Chi-square	success	<i>failure</i>	0.8300	success
RandomWalk1 J Chi-square	success	<i>failure</i>	0.9400	success
RandomWalk1 R Chi-square	success	<i>failure</i>	0.7000	success
RandomWalk1 C Chi-square	success	<i>failure</i>	0.6600	success

TABLE VII: TestU01 (Small Crush). Given the constraint of the data size and computational power of Crush and Big Crush of TestU01, we only perform Small Crush test here. Data size is 8 Gbits. The P-value of failing a test converges to 0 or 1 (eps or 1-eps). Where the test has multiple P-values, the worst case is selected.

# Our approach: randomness from laser phase noise

- Physical origin: spontaneous emissions [1, 2]



- Quantum phase change within time  $t$  can be treated as a Gaussian white noise [1,2]

$$\Delta\theta(t) \sim N(0, 2\pi t \Delta f)$$

Laser linewidth

[1] A. Yariv and P. Yeh, "Photonics: optical electronics in modern communications" (6th edition), Oxford University Press (2007).

[2] K. Petermann, "Laser diode modulation and noise", (Springer, 1988).

# Laser Intensity Noise

