**T. Eberle**<sup>1</sup>, V. Händchen<sup>1</sup>, F. Furrer<sup>2</sup>, T. Franz<sup>3</sup>, J. Duhme<sup>3</sup>, R.F. Werner<sup>3</sup>, R. Schnabel<sup>1</sup>

### Realization of

- Finite-Size Continuous-Variable
  - **Quantum Key Distribution**

### based on

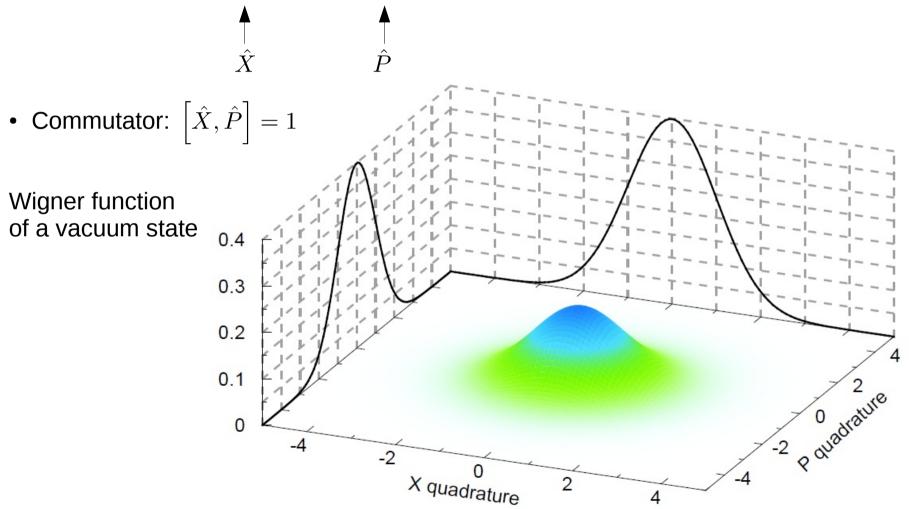
### Einstein-Podolsky-Rosen Entangled Light

- Institute for Gravitational Physics, Leibniz University Hannover and Max Planck Institute for Gravitational Physics (Albert Einstein Institute), Callinstr. 38, 30167 Hannover, Germany
- (2) Department of Physics, Graduate School of Science, University of Tokyo, 7-3-1 Hongo, Bunkyo-ku, Tokyo, Japan, 113-0033
- (3) Institute for Theoretical Physics, Leibniz University Hannover, Appelstraße 2, 30167 Hannover, Germany

QCrypt 2013 August, 5th 2013

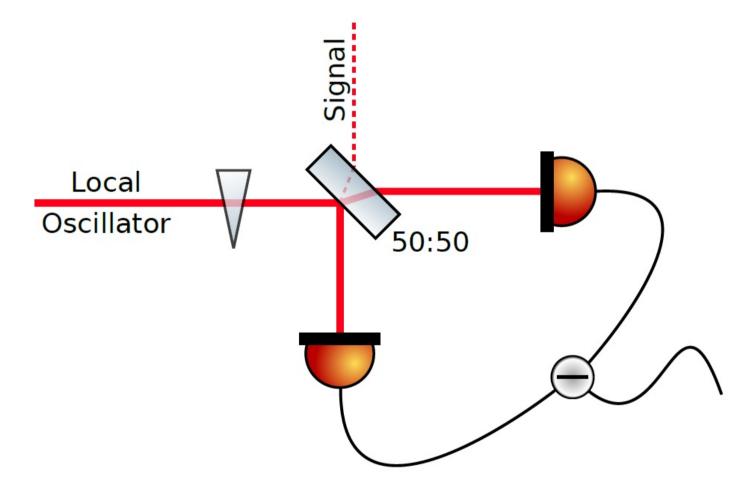


- QKD protocol requires two non-commuting observables
- **Observables:** Amplitude and Phase Quadrature of Light Fields





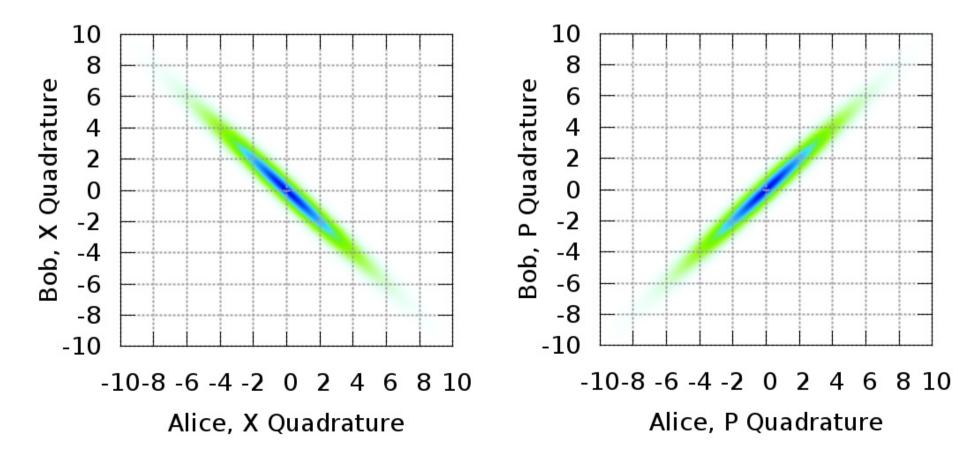
#### Homodyne detection



Phase of local oscillator with respect to signal determines measured quadrature



Simultaneous homodyne measurements at Alice and Bob



# Why entanglement based CV QKD?

### **Gaussian Modulation**

- Only standard telecommunication components Amplitude, phase modulators, PIN photo diodes
- Security analysis for collective attacks includes finite-size effects

A. Leverrier et al., Phys.I Rev. A 81, 062343 (2010).

• 80 km distance reached

P. Jouguet et al., Nature Photonics 7, 378-381 (2013).

 Compatibility with intense DWDM classical channels demonstrated

### Security analysis for **arbitrary attacks** including finite-size effects given in

F. Furrer et al., Phys. Rev. Lett. 109, 100502 (2012).

### requires quadrature entanglement

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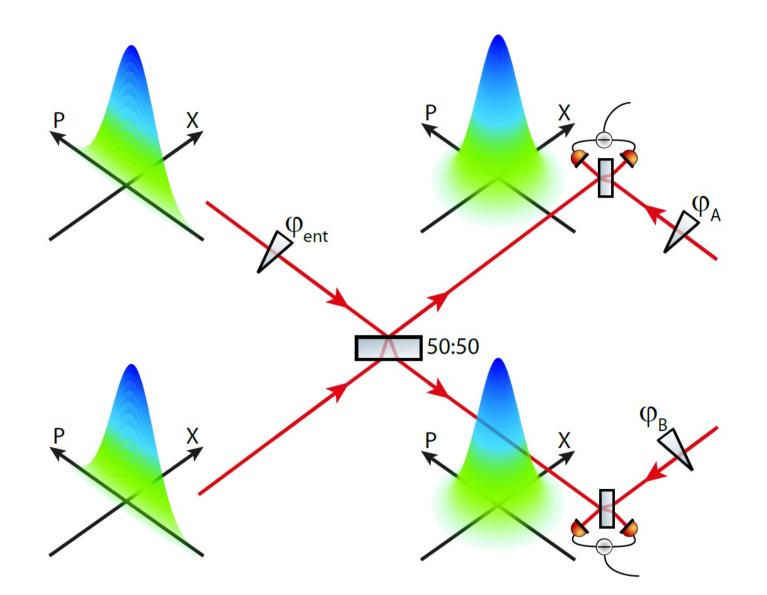


### Realization of Finite-Size Continuous-Variable Quantum Key Distribution based on Einstein-Podolsky-Rosen Entangled Light

1. Einstein-Podolsky Rosen (EPR) Entanglement

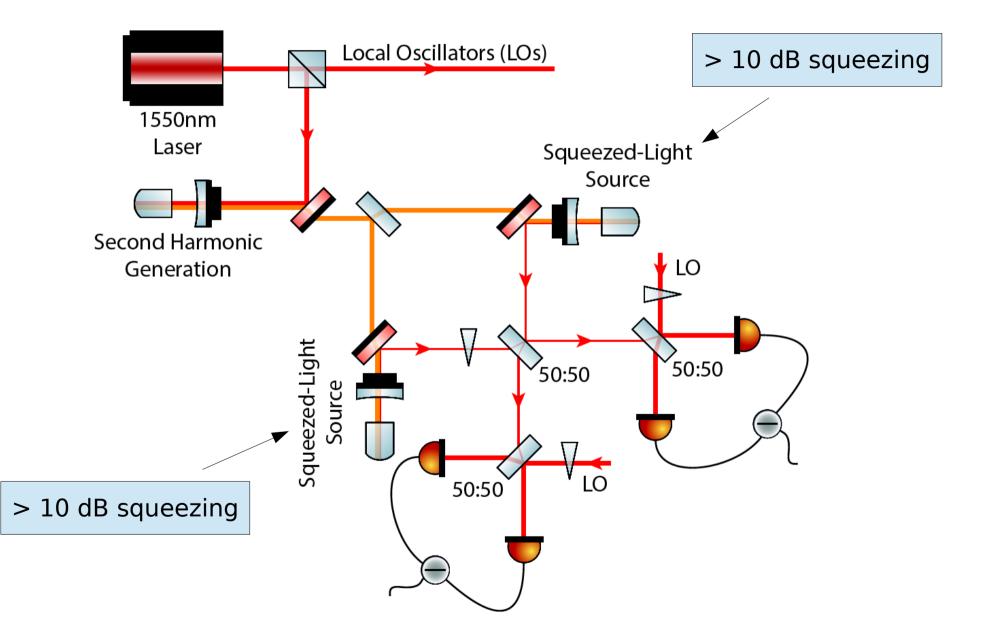
- 2. Continuous-Variable Quantum Key Distribution
  - Generate a secure key under collective attacks
  - Demonstrate the feasibility to have security under arbitrary attacks

### Generation of Quadrature Entanglement

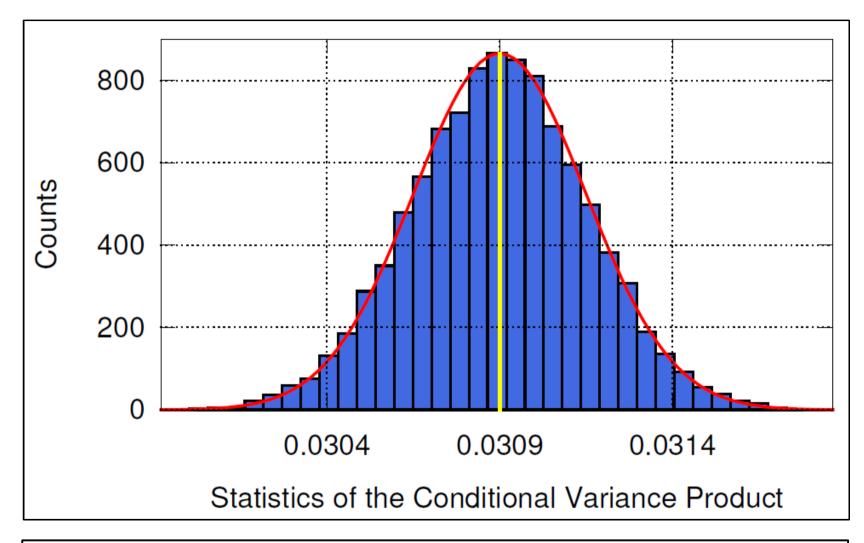


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**Experimental Setup** 

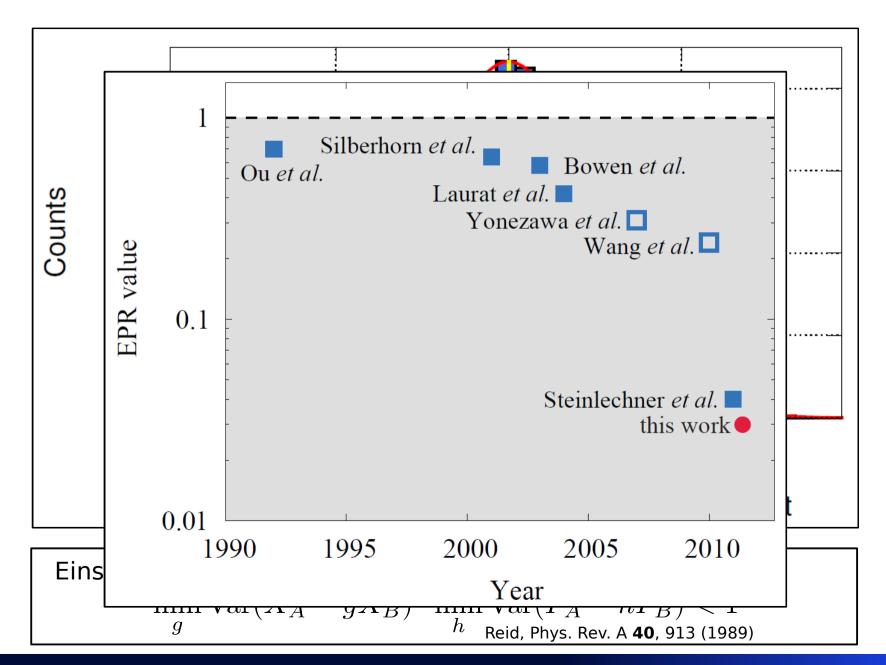


# Results: EPR Entanglement



Einstein-Podolsky-Rosen Entanglement:  $\min_{g} \operatorname{Var}(\hat{X}_A - g\hat{X}_B) \cdot \min_{h} \operatorname{Var}(\hat{P}_A - h\hat{P}_B) < 1$ <sub>Reid, Phys. Rev. A **40**, 913 (1989)</sub>





Realization of Finite-Size Continuous-Variable Quantum Key Distribution based on EPR Entanglement Tobias Eberle

#### rle

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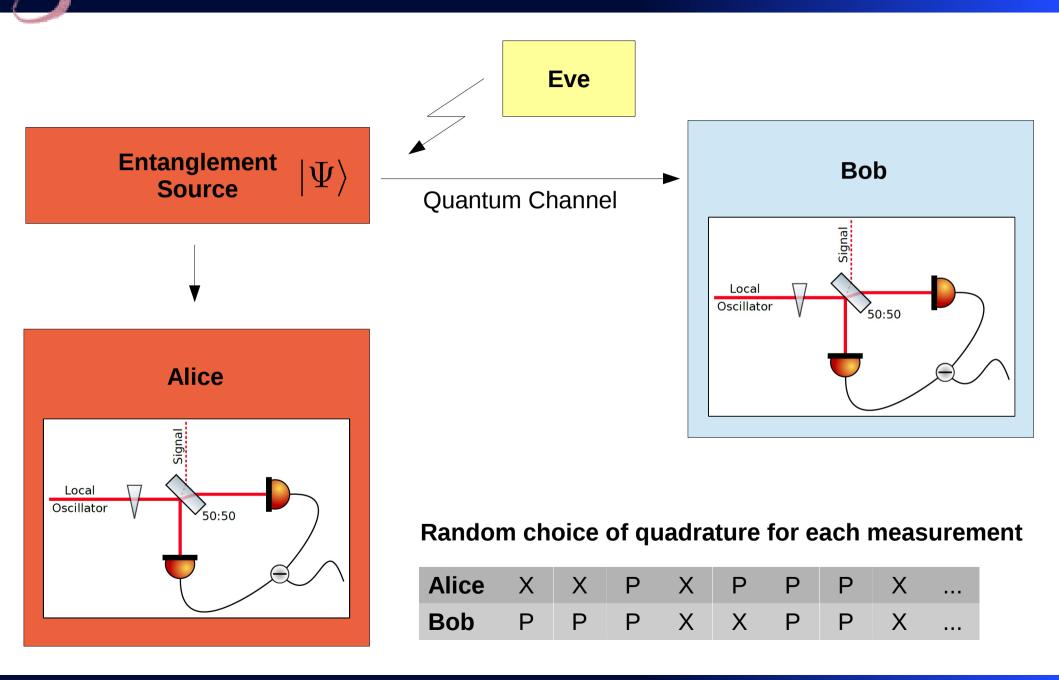
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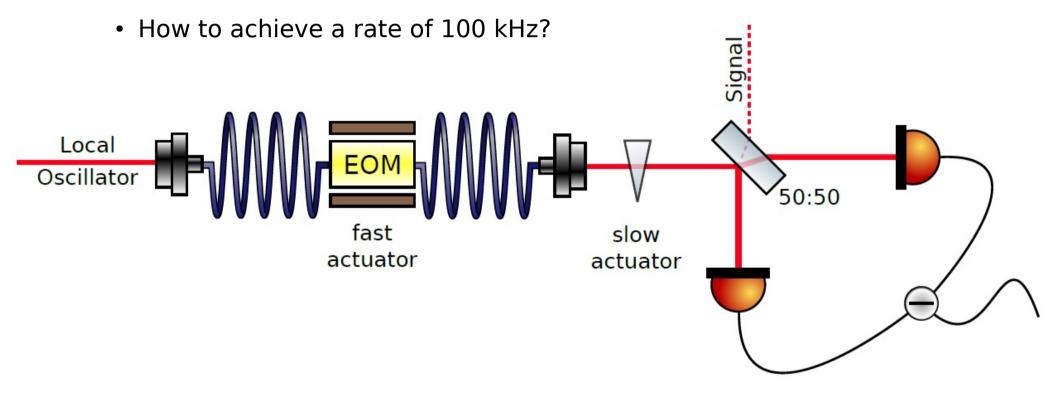
Quantum Key Distribution Scheme



Random Quadrature Choice: Implementation

Quantum part of QKD protocol: Measurement of 10<sup>8</sup>-10<sup>9</sup> samples necessary

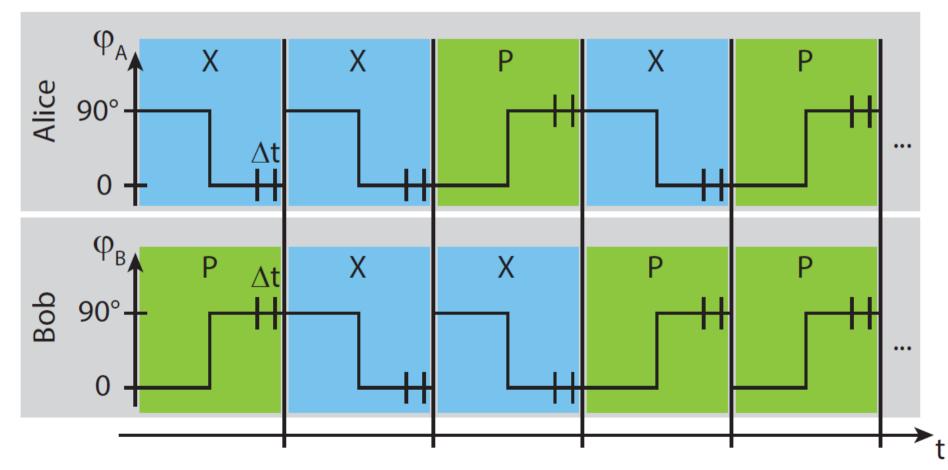
• Usual phase locks have unity gain frequencies < 1kHz



- Fast phase actuator: Shifts phase by 0 or 90° with a rate of 100 kHz
- Slow phase actuator: Compensates for drifts

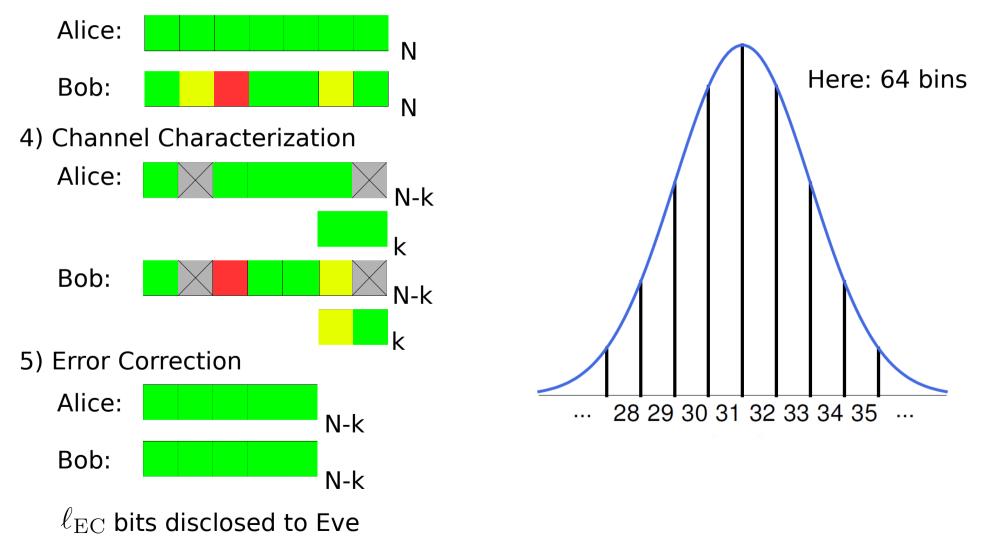
# Random Quadrature Choice: Implementation

- Low frequency lock averages over phase
- Problem: long sequences of X or P measurements
- Solution:



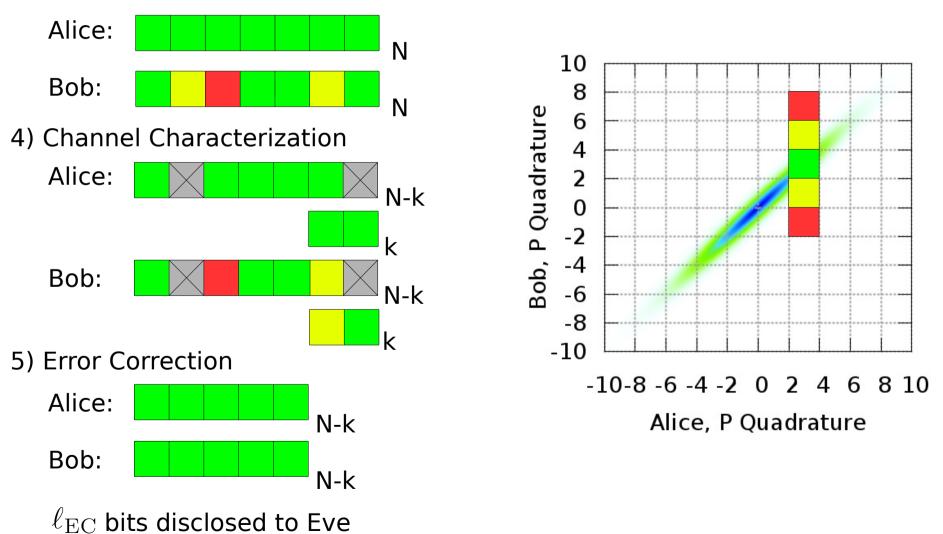
# QKD Protocol

- 2) Sifting: Discard measurements in different quadratures
- 3) Binning



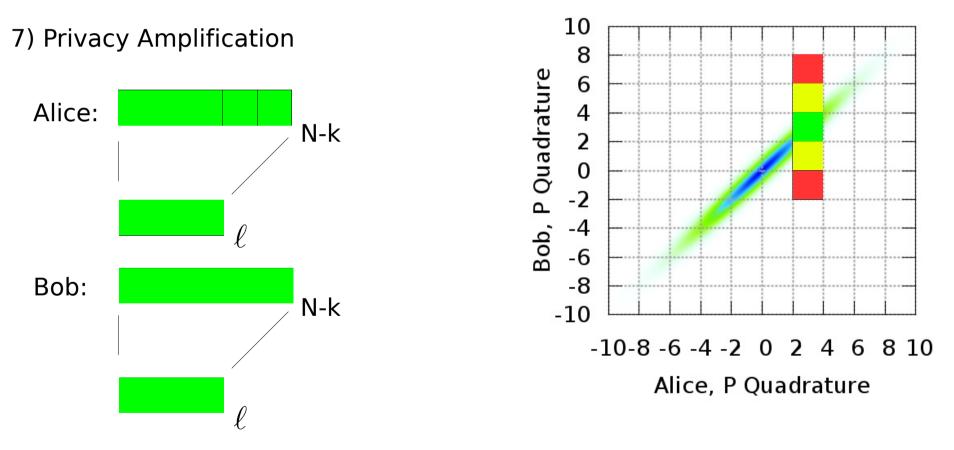
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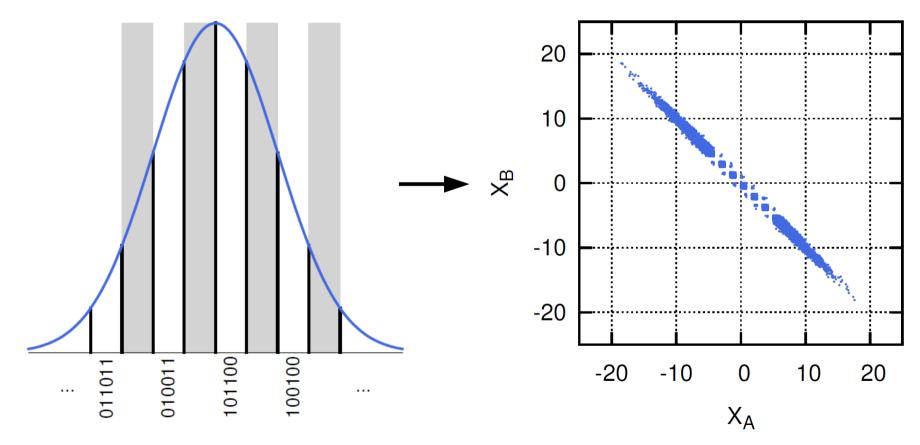


- 6) Calculate secure key length  $\ell$ 
  - Use results from channel characterization
  - Take  $\ell_{EC}$  into account



# Protocol Execution under Collective Attacks

- Measurement of 10<sup>8</sup> samples
- Bin measurement outcomes into 2<sup>6</sup>=64 intervals
- To use binary error correction a simple post selection technique was applied

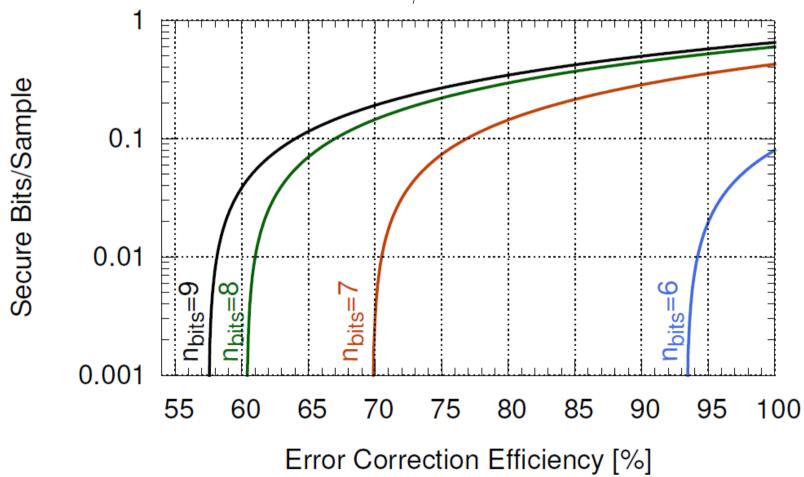


• Generation of 1.5MB key using an error correction algorithm from AIT

# Protocol Execution under Arbitrary Attacks

- Post selection not possible since it is not covered by the security proof
  - Estimation of the necessary non-binary error correction efficiency

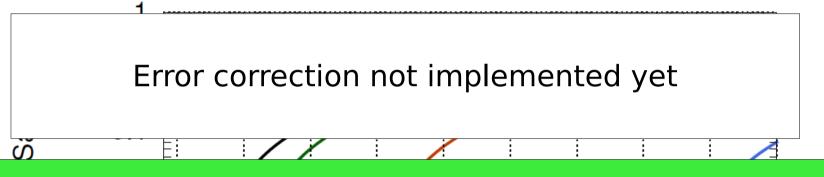
$$\ell_{\rm EC} = H(X_A) - \beta \cdot I(X_A : X_B)$$
$$\beta < 1$$



# Protocol Execution under Arbitrary Attacks

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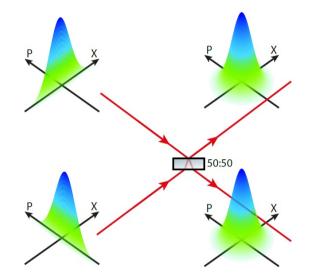
#### **Poster:**

### **CV-QKD on Hannover campus: key generation and error correction**

Jörg Duhme, Kais Abdelkhalek, Rene Schwonnek, Fabian Furrer, and Reinhard F. Werner



#### **Einstein-Podolsky-Rosen Entanglement**



 Generation of strong quadrature entanglement

**Quantum Key Distribution** 



- Generation of 1.5MB secret key under collective attacks
- Demonstrated feasibility of QKD under arbitrary attacks